

MzLib: Legacy PLT Libraries

Version 4.0

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The "mzlib" collection contains wrappers and libraries for compatibility with older versions of PLT Scheme. In many ways, the libraries of the "mzlib" collection go with the [mzscheme](#) legacy language. Newer variants of many libraries reside in the "scheme" collection.

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1 mzlib/a-signature

```
(require mzlib/a-signature)
```

Like `scheme/signature` in `#lang` form for defining a single signature within a module, but based on `mzscheme` instead of `scheme/base`.

2 mzlib/a-unit

```
(require mzlib/a-unit)
```

Like `scheme/unit` in `#lang` form for defining a single unit within a module, but based on `mzscheme` instead of `scheme/base`.

3 mzlib/async-channel

```
(require mzlib/async-channel)
```

Re-exports `scheme/async-channel`.

4 mzlib/awk

```
(require mzlib/awk)
```

```
(awk next-record-expr
  (record field-id ...)
  maybe-counter
  ((state-variable init-expr) ...)
  maybe-continue
  clause ...)

maybe-counter =
  | id

maybe-continue =
  | id

  clause = (test body ...+)
  | (test => procedure-expr)
  | (/ regexp-str / (id-or-false ...+) body ...+)
  | (range excl-start-test excl-stop-test body ...+)
  | (:range incl-start-test excl-stop-test body ...+)
  | (range: excl-start-test incl-stop-test body ...+)
  | (:range: incl-start-test incl-stop-test body ...+)
  | (else body ...+)
  | (after body ...+)

  test = integer
  | regexp-string
  | expr

excl-start-test = test

excl-stop-test = test

incl-start-test = test

incl-stop-test = test

  id-or-false = id
  | #f
```

The awk macro from Scsh [Shivers06]. In addition to awk, the Scsh-compatible procedures `match:start`, `match:end`, `match:substring`, and `regexp-exec` are defined. These

`match`: procedures must be used to extract match information in a regular expression clause when using the `=>` form.

```
(match:start rec [which]) → exact-nonnegative-integer?  
  rec : ....  
  which : exact-nonnegative-integer? = 0  
(match:end rec [which]) → exact-nonnegative-integer?  
  rec : ....  
  which : exact-nonnegative-integer? = 0  
(match:substring rec [which]) → string?  
  rec : ....  
  which : exact-nonnegative-integer? = 0
```

Extracts a start position, end position, or substring corresponding to a match. The first argument is the value supplied to the procedure after `=>` in a `awk` clause or the result of `regexp-exec`.

```
(regexp-exec re s) → (or/c .... false/c)  
  re : (or/c string? regexp?)  
  s : string?
```

Matches a `regexp` to a string, returning a record compatible with `match:start`, etc.

5 `mzlib/class`

```
(require mzlib/class)
```

Re-exports `scheme/class`, except for the contract constructors.

6 mzlib/class100

```
(require mzlib/class100)
```

The `class100` and `class100*` forms provide a syntax close to that of `class` and `class*` in PLT Scheme versions 100 through 103, but with the semantics of the current `scheme/class`-based class system. For a class defined with `class100`, keyword-based initialization arguments can be propagated to the superclass, but by-position arguments are not (i.e., the expansion of `class100` to `class` always includes an `init-rest` clause).

The `class100` form uses keywords (e.g., `public`) that are defined by the `mzlib/class` library, so typically `scheme/class` must be imported into any context that imports `mzlib/class100`.

```
(class100* superclass-expr (interface-expr ...) init-ids  
  class100-clause  
  ...)
```

```

    init-ids = id
              | (id ... id-with-default ...)
              | (id ... id-with-default ... . id)

    id-with-default = (id default-expr)

    class100-clause = (sequence expr ...)
                    | (public public-method-decl ...)
                    | (override public-method-decl ...)
                    | (augment public-method-decl ...)
                    | (pubment public-method-decl ...)
                    | (overment public-method-decl ...)
                    | (augride public-method-decl ...)
                    | (private private-method-decl ...)
                    | (private-field private-var-decl ...)
                    | (inherit inherit-method-decl ...)
                    | (rename rename-method-decl ...)

    public-method-decl = ((internal-id external-id) method-procedure)
                       | (id method-procedure)

    private-method-decl = (id method-procedure)

    private-var-decl = (id initial-value-expr)
                     | (id)
                     | id

    inherit-method-decl = id
                        | (internal-instance-id external-inherited-id)

    rename-method-decl = (internal-id external-id)

```

```

(class100 superclass-expr init-ids
 class100-clause
 ...)
```

Like *class100**, but without *interface-exprs*.

```

(class100-asi superclass instance-id-clause ...)
```

Like *class100*, but all initialization arguments are automatically passed on to the superclass initialization procedure by position.

```

(class100*-asi superclass interfaces instance-id-clause ...)
```

Like `class100*`, but all initialization arguments are automatically passed on to the super-class initialization procedure by position.

`(super-init init-arg-expr ...)`

An alias for `super-make-object`.

7 mzlib/cm

```
(require mzlib/cm)
```

Re-exports `compiler/cm`.

8 mzlib/cm-accomplice

```
(require mzlib/cm-accomplice)
```

Re-exports `compiler/cm-accomplice`.

9 mzlib/cmdline

```
(require mzlib/cmdline)
```

Provides a `command-line` from that is similar to the one in `scheme/cmdline`, but without using keywords. The `parse-command-line` procedure from `scheme/cmdline` is re-exported directly.

```
(command-line program-name-expr argv-expr clause ...)

  clause = (multi flag-spec ...)
           | (once-each flag-spec ...)
           | (once-any flag-spec ...)
           | (final flag-spec ...)
           | (help-labels string ...)
           | (args arg-formals body-expr ...)
           | (=> finish-proc-expr arg-help-expr help-proc-expr
               unknown-proc-expr)

  flag-spec = (flags id ... help-str ...+ body-expr ...)
              | (flags => handler-expr help-expr)

  flags = flag-string
          | (flag-string ...)

  arg-formals = id
               | (id ...)
               | (id ...+ . id)
```

Like `command-line` from `scheme/cmdline`, but without keywords in the syntax.

10 mzlib/cml

```
(require mzlib/cml)
```

The `mzlib/cml` library defines a number of procedures that wrap PLT Scheme concurrency procedures. The wrapper procedures have names and interfaces that more closely match those of Concurrent ML [Reppy99].

```
(spawn thunk) → thread?  
  thunk : (-> any)
```

Equivalent to `(thread/suspend-to-kill thunk)`.

```
(channel) → channel?
```

Equivalent to `(make-channel)`.

```
(channel-recv-evt ch) → evt?  
  ch : channel?
```

Equivalent to `ch`.

```
(channel-send-evt ch v) → evt?  
  ch : channel?  
  v : any/c
```

Equivalent to `(channel-put-evt ch v)`.

```
(thread-done-evt thd) → any  
  thd : thread?
```

Equivalent to `(thread-dead-evt thread)`.

```
(current-time) → real?
```

Equivalent to `(current-inexact-milliseconds)`.

```
(time-evt tm) → evt?  
  tm : real?
```

Equivalent to `(alarm-evt tm)`.

11 mzlib/compat

```
(require mzlib/compat)
```

The `mzlib/compat` library defines a number of procedures and syntactic forms that are commonly provided by other Scheme implementations. Most of the procedures are aliases for `mzscheme` procedures.

```
(=? n ...+) → boolean?  
  n : number?  
(<? n ...+) → boolean?  
  n : real?  
(>? n ...+) → boolean?  
  n : real?  
(<=? n ...+) → boolean?  
  n : real?  
(>=? n ...+) → boolean?  
  n : real?
```

Same as `=`, `<`, etc.

```
(1+ n) → number?  
  n : number?  
(1- n) → number?  
  n : number?
```

Same as `add1` and `sub1`.

```
(gentmp [base]) → symbol?  
  base : (or/c string? symbol?) = "g"
```

Same as `gensym`.

```
(flush-output-port [o]) → void?  
  o : output-port? = (current-output-port)
```

Same as `flush-output`.

```
(real-time) → exact-integer?
```

Same as `current-milliseconds`.

```
(atom? v) → any
  v : any/c
```

Same as `(not (pair? v))` (which does not actually imply an atomic value).

```
(define-structure (name-id field-id ...))
(define-structure (name-id field-id ...)
  ((init-field-id init-expr) ...))
```

Like `define-struct`, except that the *name-id* is moved inside the parenthesis for fields. In addition, *init-field-ids* can be specified with automatic initial-value expression.

The *init-field-ids* do not have corresponding arguments for the *make-name-id* constructor. Instead, each *init-field-id*'s *init-expr* is evaluated to obtain the field's value when the constructor is called. The *field-ids* are bound in *init-exprs*, but not other *init-field-ids*.

Examples:

```
> (define-structure (add left right) ([sum (+ left right)]))
> (add-sum (make-add 3 6))
9
```

```
(getprop sym property [default]) → any/c
  sym : symbol?
  property : symbol?
  default : any/c = #f
(putprop sym property value) → void?
  sym : symbol?
  property : symbol?
  value : any/c
```

The `getprop` function gets a property value associated with *sym*. The *property* argument names the property to be found. If the property is not found, *default* is returned.

The properties obtained with `getprop` are the ones installed with `putprop`.

```
(new-cafe [eval-handler]) → any
  eval-handler : (any/c . -> . any) = #f
```

Emulates Chez Scheme's `new-cafe` by installing *eval-handler* into the `current-eval` parameter while running `read-eval-print`. In addition, `current-exit` is set to escape from the call to `new-cafe`.

12 mzlib/compile

```
(require mzlib/compile)
```

```
(compile-file src [dest filter]) → path?  
src : path-string?  
dest : path-string?  
      = (let-values ([[base name dir?] (split-path src)])  
          (build-path base "compiled"  
                      (path-add-suffix name #".zo")))  
filter : (any/c . -> . any/c) = values
```

Compiles the Scheme file *src* and saves the compiled code to *dest*. If *dest* is not provided and the "compiled" subdirectory does not already exist, the subdirectory is created. The result of `compile-file` is the destination file's path.

If the *filter* procedure is provided, it is applied to each source expression, and the result is compiled.

The `compile-file` procedure is designed for compiling modules files, in that each expression in *src* is compiled independently. If *src* does not contain a single module expression, then earlier expressions can affect the compilation of later expressions when *src* is loaded directly. An appropriate *filter* can make compilation behave like evaluation, but the problem is also solved (as much as possible) by the `compile-zos` procedure.

See also `managed-compile-zo`.

13 mzlib/contract

```
(require mzlib/contract)
```

The `mzlib/list` library re-exports many bindings from `scheme/contract`:

<code></c</code>	<code>flat-rec-contract</code>
<code><=/c</code>	<code>guilty-party</code>
<code>=/c</code>	<code>integer-in</code>
<code>>/c</code>	<code>list/c</code>
<code>>=/c</code>	<code>listof</code>
<code>and/c</code>	<code>make-none/c</code>
<code>any</code>	<code>make-proj-contract</code>
<code>any/c</code>	<code>natural-number/c</code>
<code>between/c</code>	<code>none/c</code>
<code>box-immutable/c</code>	<code>not/c</code>
<code>box/c</code>	<code>one-of/c</code>
<code>build-compound-type-name</code>	<code>or/c</code>
<code>coerce-contract</code>	<code>parameter/c</code>
<code>cons/c</code>	<code>printable/c</code>
<code>contract</code>	<code>promise/c</code>
<code>contract-first-order-passes?</code>	<code>provide/contract</code>
<code>contract-violation->string</code>	<code>raise-contract-error</code>
<code>contract?</code>	<code>real-in</code>
<code>define-contract-struct</code>	<code>recursive-contract</code>
<code>define/contract</code>	<code>string/len</code>
<code>false/c</code>	<code>struct/c</code>
<code>flat-contract</code>	<code>symbols</code>
<code>flat-contract-predicate</code>	<code>syntax/c</code>
<code>flat-contract/predicate?</code>	<code>vector-immutable/c</code>
<code>flat-contract?</code>	<code>vector-immutableof</code>
<code>flat-murec-contract</code>	<code>vector/c</code>
<code>flat-named-contract</code>	<code>vectorof</code>

14 mzlib/control

```
(require mzlib/control)
```

Re-exports `scheme/control`.

15 `mzlib/date`

```
(require mzlib/date)
```

Re-exports `scheme/date`.

16 mzlib/deflate

```
(require mzlib/deflate)
```

Re-exports `file/gzip`.

17 mzlib/defmacro

```
(require mzlib/defmacro)
```

```
(define-macro id expr)
(define-macro (id . formals) body ...+)
(defmacro id formals body ...+)
```

```
formals = (id ...)
          | id
          | (id ...+ . id)
```

Defines a (non-hygienic) macro *id* through a procedure that manipulates S-expressions, as opposed to syntax objects.

In the first form, *expr* must produce a procedure. In the second form, *formals* determines the formal arguments of the procedure, as in `lambda`, and the *exprs* are the procedure body. The last form, with `defmacro`, is like the second form, but with slightly different parentheses.

In all cases, the procedure is generated in the transformer environment, not the normal environment.

In a use of the macro,

```
(id datum ...)
```

`syntax->datum` is applied to the expression, and the transformer procedure is applied to the `cdr` of the resulting list. If the number of `datums` does not match the procedure's arity, or if *id* is used in a context that does not match the above pattern, then a syntax error is reported.

After the macro procedure returns, the result is compared to the procedure's arguments. For each value that appears exactly once within the arguments (or, more precisely, within the S-expression derived from the original source syntax), if the same value appears in the result, it is replaced with a syntax object from the original expression. This heuristic substitution preserves source location information in many cases, despite the macro procedure's operation on raw S-expressions.

After substituting syntax objects for preserved values, the entire macro result is converted to syntax with `datum->syntax`. The original expression supplies the lexical context and source location for converted elements.

Important: Although `define-macro` is non-hygienic, it is still restricted by PLT Scheme's phase separation rules. This means that a macro cannot access run-time bindings, because it is executed in the syntax-expansion phase. Translating code that involves `define-macro` or `defmacro` from an implementation without this restriction usually implies separating macro

related functionality into a `begin-for-syntax` or a module (that will be imported with `require-for-syntax`) and properly distinguishing syntactic information from run-time information.

18 mzlib/etc

```
(require mzlib/etc)
```

The `mzlib/etc` library re-exports the following `scheme/base` and other libraries:

```
boolean=?  
true  
false  
build-list  
build-string  
build-vector  
compose  
local  
symbol=?
```

```
(begin-lifted expr ...+)
```

Lifts the `exprs` so that they are evaluated once at the “top level” of the current context, and the result of the last `expr` is used for every evaluation of the `begin-lifted` form.

When this form is used as a run-time expression within a module, the “top level” corresponds to the module’s top level, so that each `expr` is evaluated once for each invocation of the module. When it is used as a run-time expression outside of a module, the “top level” corresponds to the true top level. When this form is used in a `define-syntax`, `letrec-syntax`, etc. binding, the “top level” corresponds to the beginning of the binding’s right-hand side. Other forms may redefine “top level” (using `local-expand/capture-lifts`) for the expressions that they enclose.

```
(begin-with-definitions defn-or-expr ...)
```

Supports a mixture of expressions and mutually recursive definitions, much like a module body. Unlike in a module, however, syntax definitions cannot be used to generate other immediate definitions (though they can be used for expressions).

The result of the `begin-with-definitions` form is the result of the last `defn-or-expr` if it is an expression, `#<void>` otherwise. If no `defn-or-expr` is provided (after flattening `begin` forms), the result is `#<void>`.

```
(define-syntax-set (id ...) defn ...)
```

Similar to `define-syntaxes`, but instead of a single body expression, a sequence of definitions follows the sequence of defined identifiers. For each `identifier`, the `defns` should

include a definition for *id/proc*. The value for *id/proc* is used as the (expansion-time) value for *id*.

The `define-syntax-set` form is useful for defining a set of syntax transformers that share helper functions, though `begin-for-syntax` now serves essentially the same purposes.

Examples:

```
(define-syntax-set (let-current-continuation
                  let-current-escape-continuation)
  (define (mk call-id)
    (lambda (stx)
      (syntax-case stx ()
        [(_ id body1 body ...)
         (with-syntax ([call call-id])
           #'(call (lambda (id) body1 body ...))))]))
  (define let-current-continuation/proc
    (mk (quote-syntax call/cc)))
  (define let-current-escape-continuation/proc
    (mk (quote-syntax call/ec))))
```

```
(evcase key-expr (value-expr body-expr ...) ...+)
(evcase key-expr (value-expr body-expr ...) ... [else body-expr ...])
```

The `evcase` form is similar to `case`, except that expressions are provided in each clause instead of a sequence of data. After *key-expr* is evaluated, each *value-expr* is evaluated until a value is found that is `eqv?` to the key value; when a matching value is found, the corresponding *body-exprs* are evaluated and the value(s) for the last is the result of the entire `evcase` expression.

The `else` literal is recognized either as unbound (like in the `mzscheme` language) or bound as `else` from `scheme/base`.

```
(identity v) → any/c
  v : any/c
```

Returns *v*.

```
(let+ clause body-expr ...)
```

```

clause = (val target expr)
         | (rec target expr)
         | (vals (target ...) expr)
         | (recs (target expr) ...)
         | (_ expr ...)

target = id
         | (values id ...)

```

A binding construct that specifies scoping on a per-binding basis instead of a per-expression basis. It helps eliminate rightward-drift in programs. It looks similar to `let`, except each clause has an additional keyword tag before the binding variables.

Each *clause* has one of the following forms:

- `(val target expr)` : Binds *target* non-recursively to *expr*.
- `(rec target expr)` : Binds *target* recursively to *expr*.
- `(vals (target expr) ...)` : The *targets* are bound to the *exprs*. The environment of the *exprs* is the environment active before this clause.
- `(recs (target expr) ...)` : The *targetss* are bound to the *exprs*. The environment of the *exprs* includes all of the *targetss*.
- `(_ expr ...)` : Evaluates the *exprs* without binding any variables.

The clauses bind left-to-right. When a *target* is `(values id ...)`, multiple values returned by the corresponding expression are bound to the multiple variables.

Examples:

```

> (let+ ([val (values x y) (values 1 2)])
      (list x y))
(1 2)
> (let ([x 1])
      (let+ ([val x 3]
             [val y x]
             y)))
3

```

```

(loop-until start done? next f) → void?
start : any/c
done? : (any/c . -> . any)
next : (any/c . -> . any/c)
f : (any/c . -> . any)

```

Repeatedly invokes the *f* procedure until the *done?* procedure returns *#t*:

```
(define (loop-until start done? next f)
  (let loop ([i start])
    (unless (done? i)
      (f i)
      (loop (next i)))))
```

```
(namespace-defined? sym) → boolean?
sym : symbol?
```

Returns *#t* if *namespace-variable-value* would return a value for *sym*, *#f* otherwise.

```
(nand expr ...)
```

Same as `(not (and expr ...))`.

```
(nor expr ...)
```

Same as `(not (or expr ...))`.

```
(opt-lambda formals body ...+)
```

Supports optional (but not keyword) arguments like `lambda` from `scheme/base`.

```
(recur id bindings body ...+)
```

Equivalent to `(let id bindings body ...+)`.

```
(rec id value-expr)
(rec (id arg-id ...) expr)
(rec (id arg-id ... . rest-id) expr)
```

Equivalent, respectively, to

```
(letrec ([id value-expr] id)
  (letrec ([id (lambda (arg-id ...) value-expr)] id)
    (letrec ([id (lambda (arg-id ... . rest-id) value-expr)] id)
```

```
(this-expression-source-directory)
(this-expression-source-directory datum)
```

Expands to an expression that evaluates to the directory of the file containing the source *datum*. If *datum* is not supplied, then the entire `(this-expression-source-directory)` expression is used as *datum*.

If *datum* has a source module, then the expansion attempts to determine the module's runtime location. This location is determined by preserving the lexical context of *datum* in a syntax object, extracting its source module path at run time, and then resolving the module path.

Otherwise, *datum*'s source file is determined through source location information associated with *datum*, if it is present. As a last resort, `current-load-relative-directory` is used if it is not `#f`, and `current-directory` is used if all else fails.

A directory path derived from source location is always stored in bytes in the expanded code, unless the file is within the result of `find-collects-dir`, in which case the expansion records the path relative to `(find-collects-dir)` and then reconstructs it using `(find-collects-dir)` at run time.

```
(this-expression-file-name)
(this-expression-file-name datum)
```

Similar to `this-expression-source-directory`, except that only source information associated with *datum* or `(this-expression-file-name)` is used to extract a filename. If no filename is available, the result is `#f`.

```
(hash-table (quote flag) ... (key-expr val-expr) ...)
```

Creates a new hash-table providing the quoted flags (if any) to `make-hash-table`, and then mapping each key to the corresponding values.

See `scheme/runtime-path` for a definition form that works better when creating executables.

19 mzlib/file

```
(require mzlib/file)
```

The `mzlib/file` library mostly re-exports from `scheme/file`:

```
find-relative-path
explode-path
normalize-path
filename-extension
file-name-from-path
path-only
delete-directory/files
copy-directory/files
make-directory*
make-temporary-file
get-preference
put-preferences
fold-files
find-files
pathlist-closure
```

```
(call-with-input-file* file proc [mode]) → any
  file : path-string?
  proc : (input-port? → any)
  mode : (one-of/c 'text 'binary) = 'binary
(call-with-output-file* file
  proc
  [mode
  exists]) → any
  file : path-string?
  proc : (output-port? → any)
  mode : (one-of/c 'text 'binary) = 'binary
  exists : (one-of/c 'error 'append 'update
    'replace 'truncate 'truncate/replace)
    = 'error
```

Like `call-with-input-file` and `call-with-output-file`, except that the opened port is closed if control escapes from the body of `proc`.

```
(build-relative-path base sub ...) → (and/c path? relative-path?)
  base : (or/c path-string?
    (one-of/c 'up 'same))
```

```

sub : (or/c (and/c path-string?
              relative-path?)
            (one-of/c 'up 'same))
(build-absolute-path base sub ...) → (and/c path? absolute-path?)
base : (or/c (and/c path-string?
              (not/c relative-path?))
          (one-of/c 'up 'same))
sub : (or/c (and/c path-string?
              (not/c complete-path?))
          (one-of/c 'up 'same))

```

Like `build-path`, but with extra constraints to ensure a relative or absolute result.

20 mzlib/for

```
(require mzlib/for)
```

The `mzlib/for` library re-exports from `scheme/base`:

```
for/fold for*/fold  
for for*  
for/list for*/list  
for/lists for*/lists  
for/and for*/and  
for/or for*/or  
for/first for*/first  
for/last for*/last
```

```
for/fold/derived for*/fold/derived
```

```
in-range  
in-naturals  
in-list  
in-vector  
in-string  
in-bytes  
in-input-port-bytes  
in-input-port-chars  
in-hash-table  
in-hash-table-keys  
in-hash-table-values  
in-hash-table-pairs
```

```
in-parallel  
stop-before  
stop-after  
in-indexed
```

```
sequence?  
sequence-generate
```

```
define-sequence-syntax  
make-do-sequence  
:do-in
```

21 `mzlib/foreign`

`(require mzlib/foreign)`

Re-exports `scheme/foreign`.

22 mzlib/include

```
(require mzlib/include)
```

Similar to `scheme/include`, but with a different syntax for paths.

```
(include path-spec)
```

```
path-spec = string
           | (build-path elem ...+)
           | (lib file-string collection-string ...)
```

```
   elem = string
         | up
         | same
```

Inlines the syntax in the designated file in place of the `include` expression. The *path-spec* can be any of the following:

- A literal string that specifies a path to include, parsed according to the platform's conventions (which means that it is not portable).
- A path construction of the form `(build-path elem ...+)`, where `build-path` is `module-identifier=?` either to the `build-path` export from `mzscheme` or to the top-level `build-path`, and where each *elem* is a path string, `up` (unquoted), or `same` (unquoted). The *elems* are combined in the same way as for the `build-path` function to obtain the path to include.
- A path construction of the form `(lib file-string collection-string ...)`, where `lib` is free or refers to a top-level `lib` variable. The *collection-strings* are passed to `collection-path` to obtain a directory; if no *collection-strings* are supplied, `"mzlib"` is used. The *file-string* is then appended to the directory using `build-path` to obtain the path to include.

If *path-spec* specifies a relative path to include, the path is resolved relative to the source for the `include` expression, if that source is a complete path string. If the source is not a complete path string, then *path-spec* is resolved relative to the current load relative directory if one is available, or to the current directory otherwise.

The included syntax is given the lexical context of the `include` expression.

```
(include-at/relative-to context source path-spec)
(include-at/relative-to/reader context source path-spec reader-expr)
(include/reader path-spec reader-expr)
```

Variants of `include` analogous to the variants of `scheme/include`.

23 mzlib/inflate

```
(require mzlib/inflate)
```

Re-exports `file/gunzip`.

24 mzlib/integer-set

```
(require mzlib/integer-set)
```

The `mzlib/integer-set` library provides functions for working with finite sets of integers. This module is designed for sets that are compactly represented as groups of intervals, even when their cardinality is large. For example, the set of integers from -1000000 to 1000000 except for 0, can be represented as `{[-1000000, -1], [1, 1000000]}`. This data structure would not be a good choice for the set of all odd integers between 0 and 1000000, which would be `{[1, 1], [3, 3], ... [999999, 999999]}`.

In addition to the *integer set* abstract type, a *well-formed set* is a list of pairs of exact integers, where each pair represents a closed range of integers, and the entire set is the union of the ranges. The ranges must be disjoint and increasing. Further, adjacent ranges must have at least one integer between them. For example: `'((-1 . 2) (4 . 10))` is a well-formed-set as is `'((1 . 1) (3 . 3))`, but `'((1 . 5) (6 . 7))`, `'((1 . 5) (-3 . -1))`, `'((5 . 1))`, and `'((1 . 5) (3 . 6))` are not.

```
(make-integer-set wfs) → integer-set?  
wfs : well-formed-set?
```

Creates an integer set from a well-formed set.

```
(integer-set-contents s) → well-formed-set?  
s : integer-set?
```

Produces a well-formed set from an integer set.

```
(set-integer-set-contents! s wfs) → void?  
s : integer-set?  
wfs : well-formed-set?
```

Mutates an integer set.

```
(integer-set? v) → boolean?  
v : any/c
```

Returns `#t` if `v` is an integer set, `#f` otherwise.

```
(well-formed-set? v) → boolean?  
v : any/c
```

Returns `#t` if `v` is a well-formed set, `#f` otherwise.

```
(make-range) → integer-set?  
(make-range elem) → integer-set?  
  elem : exact-integer?  
(make-range start end) → integer-set?  
  start : exact-integer?  
  end : exact-integer?
```

Produces, respectively, an empty integer set, an integer set containing only `elem`, or an integer set containing the integers from `start` to `end` inclusive, where (`<= start end`).

```
(intersect x y) → integer-set?  
  x : integer-set?  
  y : integer-set?
```

Returns the intersection of the given sets.

```
(difference x y) → integer-set?  
  x : integer-set?  
  y : integer-set?
```

Returns the difference of the given sets (i.e., elements in `x` that are not in `y`).

```
(union x y) → integer-set?  
  x : integer-set?  
  y : integer-set?
```

Returns the union of the given sets.

```
(split x y) → integer-set?  
  x : integer-set?  
  y : integer-set?
```

Produces three values: the first is the intersection of `x` and `y`, the second is the difference `x` remove `y`, and the third is the difference `y` remove `x`.

```
(complement s start end) → any  
  s : integer-set?  
  start : exact-integer?  
  end : exact-integer?
```

Returns the a set containing the elements between `start` to `end` inclusive that are not in `s`, where `(<= start-k end-k).`}

```
(xor x y) → integer-set?  
  x : integer-set?  
  y : integer-set?
```

Returns an integer set containing every member of `x` and `y` that is not in both sets.

```
(member? k s) → boolean?  
  k : exact-integer?  
  s : integer-set?
```

Returns `#t` if `k` is in `s`, `#f` otherwise.

```
(get-integer integer-set) → (or/c exact-integer? false/c)  
  integer-set : any/c
```

Returns a member of `integer-set`, or `#f` if `integer-set` is empty.

```
(foldr proc base-v s) → any/c  
  proc : (exact-integer? any/c . -> . any/c)  
  base-v : any/c  
  s : integer-set?
```

Applies `proc` to each member of `s` in ascending order, where the first argument to `proc` is the set member, and the second argument is the fold result starting with `base-v`. For example, `(foldr cons null s)` returns a list of all the integers in `s`, sorted in increasing order.

```
(partition s) → (listof integer-set?)  
  s : integer-set-list?
```

Returns the coarsest refinement of the sets in `s` such that the sets in the result list are pairwise disjoint. For example, partitioning the sets that represent `'((1 . 2) (5 . 10))` and `'((2 . 2) (6 . 6) (12 . 12))` produces the a list containing the sets for `'((1 . 1) (5 . 5) (7 . 10))`, `'((2 . 2) (6 . 6))`, and `'((12 . 12))`.

```
(card s) → exact-nonnegative-integer?  
  s : integer-set?
```

Returns the number of integers in the given integer set.

```
(subset? x y) → boolean?  
x : integer-set?  
y : integer-set?
```

Returns true if every integer in `x` is also in `y`, otherwise `#f`.

25 mzlib/kw

```
(require mzlib/kw)

(lambda/kw kw-formals body ...+)
(define/kw (head args) body ...+)

kw-formals = id
             | (id ... [#:optional optional-spec ...]
               [#:key key-spec ...]
               [rest/mode-spec ...])
             | (id ... . id)

optional-spec = id
                | (id default-expr)

key-spec = id
            | (id default-expr)
            | (id keyword default-expr)

rest/mode-spec = #:rest id
                  | #:other-keys id
                  | #:other-keys+body id
                  | #:all-keys id
                  | #:body kw-formals
                  | #:allow-other-keys
                  | #:forbid-other-keys
                  | #:allow-duplicate-keys
                  | #:forbid-duplicate-keys
                  | #:allow-body
                  | #:forbid-body
                  | #:allow-anything
                  | #:forbid-anything

head = id
        | (head . kw-formals)
```

The `lambda` and procedure-application forms of `scheme/base` support keyword arguments, and it is *not* compatible with the `mzlib/kw` library.

Like `lambda`, but with optional and keyword-based argument processing. This form is similar to an extended version of Common Lisp procedure arguments (but note the differences below). When used with plain variable names, `lambda/kw` expands to a plain `lambda`, so `lambda/kw` is suitable for a language module that will use it to replace `lambda`. Also, when used with only optionals, the resulting procedure is similar to `opt-lambda` (but a bit faster).

In addition to `lambda/kw`, `define/kw` is similar to `define`, except that the `formals` are

as in `lambda/kw`. Like `define`, this form can be used with nested parenthesis for curried functions (the MIT-style generalization of `define`).

The syntax of `lambda/kw` is the same as `lambda`, except for the list of formal argument specifications. These specifications can hold (zero or more) plain argument names, then an optionals (and defaults) section that begins after an `#:optional` marker, then a keyword section that is marked by `#:keyword`, and finally a section holding rest and “rest”-like arguments which are described below, together with argument processing flag directives. Each section is optional, but the order of the sections must be as listed. Of course, all binding *ids* must be unique.

The following sections describe each part of the *kw-formals*.

25.1 Required Arguments

Required arguments correspond to *ids* that appear before any keyword marker in the argument list. They determine the minimum arity of the resulting procedure.

25.2 Optional Arguments

The optional-arguments section follows an `#:optional` marker in the *kw-formals*. Each optional argument can take the form of a parenthesized variable and a default expression; the latter is used if a value is not given at the call site. The default expression can be omitted (along with the parentheses), in which case `#f` is the default.

The default expression’s environment includes all previous arguments, both required and optional names. With k optionals after n required arguments, and with no keyword arguments or rest-like arguments, the resulting procedure accept between n and $n+k$ arguments, inclusive.

The treatment of optionals is efficient, with an important caveat: default expressions appear multiple times in the resulting `case-lambda`. For example, the default expression for the last optional argument appears $k-1$ times (but no expression is ever evaluated more than once in a procedure call). This expansion risks exponential blow-up if `lambda/kw` is used in a default expression of a `lambda/kw`, etc. The bottom line, however, is that `lambda/kw` is a sensible choice, due to its enhanced efficiency, even when you need only optional arguments.

Using both optional and keyword arguments is possible, but note that the resulting behavior differs from traditional keyword facilities (including the one in Common Lisp). See the following section for details.

25.3 Keyword Arguments

A keyword argument section is marked by a `#:key`. If it is used with optional arguments, then the keyword specifications must follow the optional arguments (which mirrors the use in call sites; where optionals are given before keywords).

When a procedure accepts both optional and keyword arguments, the argument-handling convention is slightly different than in traditional keyword-argument facilities: a keyword after required arguments marks the beginning of keyword arguments, no matter how many optional arguments have been provided before the keyword. This convention restricts the procedure's non-keyword optional arguments to non-keyword values, but it also avoids confusion when mixing optional arguments and keywords. For example, when a procedure that takes two optional arguments and a keyword argument `#:x` is called with `#:x 1`, then the optional arguments get their default values and the keyword argument is bound to `1`. (The traditional behavior would bind `#:x` and `1` to the two optional arguments.) When the same procedure is called with `1 #:x 2`, the first optional argument is bound to `1`, the second optional argument is bound to its default, and the keyword argument is bound to `2`. (The traditional behavior would report an error, because `2` is provided where `#:x` is expected.)

Like optional arguments, each keyword argument is specified as a parenthesized variable name and a default expression. The default expression can be omitted (with the parentheses), in which case `#f` is the default value. The keyword used at a call site for the corresponding variable has the same name as the variable; a third form of keyword arguments has three parts—a variable name, a keyword, and a default expression—to allow the name of the locally bound variable to differ from the keyword used at call sites.

When calling a procedure with keyword arguments, the required argument (and all optional arguments, if specified) must be followed by an even number of arguments, where the first argument is a keyword that determines which variable should get the following value, etc. If the same keyword appears multiple times (and if multiple instances of the keyword are allowed; see §25.6 “Mode Keywords”), the value after the first occurrence is used for the variable:

Examples:

```
> ((lambda/kw (#:key x [y 2] [z #:zz 3] #:allow-duplicate-keys)
    (list x y z))
   #:x 'x #:zz 'z #:x "foo")
(x 2 z)
```

Default expressions are evaluated only for keyword arguments that do not receive a value for a particular call. Like optional arguments, each default expression is evaluated in an environment that includes all previous bindings (required, optional, and keywords that were specified on its left).

See §25.6 “Mode Keywords” for information on when duplicate or unknown keywords are allowed at a call site.

25.4 Rest and Rest-like Arguments

The last *kw-formals* section—after the required, optional, and keyword arguments—may contain specifications for rest-like arguments and/or mode keywords. Up to five rest-like arguments can be declared, each with an *id* to bind:

- `#:rest` — The variable is bound to the list of “rest” arguments, which is the list of all values after the required and the optional values. This list includes all keyword-value pairs, exactly as they are specified at the call site.
Scheme’s usual dot-notation is accepted in *kw-formals* only if no other meta-keywords are specified, since it is not clear whether it should specify the same binding as a `#:rest` or as a `#:body`. The dot notation is allowed without meta-keywords to make the `lambda/kw` syntax compatible with `lambda`.
- `#:body` — The variable is bound to all arguments after keyword–value pairs. (This is different from Common Lisp’s `&body`, which is a synonym for `&rest`.) More generally, a `#:body` specification can be followed by another *kw-formals*, not just a single *id*; see §25.5 “Body Argument” for more information.
- `#:all-keys` — the variable is bound to the list of all keyword-values from the call site, which is always a proper prefix of a `#:rest` argument. (If no `#:body` arguments are declared, then `#:all-keys` binds the same as `#:rest`.) See also `keyword-get`.
- `#:other-keys` — The variable is bound like an `#:all-keys` variable, except that all keywords specified in the *kw-formals* are removed from the list. When a keyword is used multiple times at a call site (and this is allowed), only the first instances is removed for the `#:other-keys` binding.
- `#:other-keys+body` — the variable is bound like a `#:rest` variable, except that all keywords specified in the *kw-formals* are removed from the list. When a keyword is used multiple times at a call site (and this is allowed), only the first instance is removed for the `#:other-keys+body` binding. (When no `#:body` variables are specified, then `#:other-keys+body` is the same as `#:other-keys`.)

In the following example, all rest-like arguments are used and have different bindings:

Examples:

```
> ((lambda/kw (#:key x y
               #:rest r
               #:other-keys+body rk
               #:all-keys ak
               #:other-keys ok
               #:body b)
     (list r rk b ak ok))
   #:z 1 #:x 2 2 3 4)
((#:z 1 #:x 2 2 3 4) (#:z 1 2 3 4) (2 3 4) (#:z 1 #:x 2) (#:z 1))
```

Note that the following invariants always hold:

- `rest = (append all-keys body)`
- `other-keys+body = (append other-keys body)`

To write a procedure that uses a few keyword argument values, and that also calls another procedure with the same list of arguments (including all keywords), use `#:other-keys` (or `#:other-keys+body`). The Common Lisp approach is to specify `:allow-other-keys`, so that the second procedure call will not cause an error due to unknown keywords, but the `:allow-other-keys` approach risks confusing the two layers of keywords.

25.5 Body Argument

The most notable divergence from Common Lisp in `lambda/kw` is the `#:body` argument, and the fact that it is possible at a call site to pass plain values after the keyword-value pairs. The `#:body` binding is useful for procedure calls that use keyword-value pairs as sort of an attribute list before the actual arguments to the procedure. For example, consider a procedure that accepts any number of numeric arguments and will apply a procedure to them, but the procedure can be specified as an optional keyword argument. It is easily implemented with a `#:body` argument:

Examples:

```
> (define/kw (mathop #:key [op +] #:body b)
  (apply op b))
> (mathop 1 2 3)
6
> (mathop #:op max 1 2 3)
3
```

(Note that the first body value cannot itself be a keyword.)

A `#:body` declaration works as an arbitrary `kw-formals`, not just a single variable like `b` in the above example. For example, to make the above `mathop` work only on three arguments that follow the keyword, use `(x y z)` instead of `b`:

Examples:

```
> (define/kw (mathop #:key [op +] #:body (x y z))
  (op x y z))
```

In general, `#:body` handling is compiled to a sub procedure using `lambda/kw`, so that a procedure can use more than one level of keyword arguments. For example:

Examples:

```

> (define/kw (mathop #:key [op +]
                   #:body (x y z #:key [convert values])))
      (op (convert x) (convert y) (convert z)))
> (mathop #:op * 2 4 6 #:convert exact->inexact)
48.0

```

Obviously, nested keyword arguments works only when non-keyword arguments separate the sets.

Run-time errors during such calls report a mismatch for a procedure with a name that is based on the original name plus a `~body` suffix:

Examples:

```

> (mathop #:op * 2 4)
procedure mathop~body: expects at least 3 arguments, given
2: 2 4

```

25.6 Mode Keywords

Finally, the argument list of a `lambda/kw` can contain keywords that serve as mode flags to control error reporting.

- `#:allow-other-keys` — The keyword-value sequence at the call site *can* include keywords that are not listed in the keyword part of the `lambda/kw` form.
- `#:forbid-other-keys` — The keyword-value sequence at the call site *cannot* include keywords that are not listed in the keyword part of the `lambda/kw` form, otherwise the `exn:fail:contract` exception is raised.
- `#:allow-duplicate-keys` — The keyword-value list at the call site *can* include duplicate values associated with same keyword, the first one is used.
- `#:forbid-duplicate-keys` — The keyword-value list at the call site *cannot* include duplicate values for keywords, otherwise the `exn:fail:contract` exception is raised. This restriction applies only to keywords that are listed in the keyword part of the `lambda/kw` form — if other keys are allowed, this restriction does not apply to them.
- `#:allow-body` — Body arguments *can* be specified at the call site after all keyword-value pairs.
- `#:forbid-body` — Body arguments *cannot* be specified at the call site after all keyword-value pairs.
- `#:allow-anything` — Allows all of the above, and treat a single keyword at the end of an argument list as a `#:body`, a situation that is usually an error. When this is used

and no rest-like arguments are used except `#:rest`, an extra loop is saved and calling the procedures is faster (around 20%).

- `#:forbid-anything` — Forbids all of the above, ensuring that calls are as restricted as possible.

These above mode markers are rarely needed, because the default modes are determined by the declared rest-like arguments:

- The default is to allow other keys if a `#:rest`, `#:other-keys+body`, `#:all-keys`, or `#:other-keys` variable is declared (and an `#:other-keys` declaration requires allowing other keys).
- The default is to allow duplicate keys if a `#:rest` or `#:all-keys` variable is declared.
- The default is to allow body arguments if a `#:rest`, `#:body`, or `#:other-keys+body` variable is declared (and a `#:body` argument requires allowing them).

Here's an alternate specification, which maps rest-like arguments to the behavior that they imply:

- `#:rest`: Everything is allowed (a body, other keys, and duplicate keys);
- `#:other-keys+body`: Other keys and body are allowed, but duplicates are not;
- `#:all-keys`: Other keys and duplicate keys are allowed, but a body is not;
- `#:other-keys`: Other keys must be allowed (on by default, cannot use with `#:forbid-other-keys`), and duplicate keys and body are not allowed;
- `#:body`: Body must be allowed (on by default, cannot use with `#:forbid-body`) and other keys and duplicate keys and body are not allowed;
- Except for the previous two “must”s, defaults can be overridden by an explicit `#:allow-...` or a `#:forbid-...` mode.

25.7 Property Lists

```
(keyword-get args kw not-found) → any
args : (listof (cons/c keyword? any/c))
kw : keyword?
not-found : (-> any)
```

Searches a list of keyword arguments (a “property list” or “plist” in Lisp jargon) for the given keyword, and returns the associated value. It is the facility that is used by `lambda/kw` to search for keyword values.

The *args* list is scanned from left to right, if the keyword is found, then the next value is returned. If the *kw* was not found, then the *not-found* thunk is used to produce a value by applying it. If the *kw* was not found, and *not-found* thunk is not given, `#f` is returned. (No exception is raised if the *args* list is imbalanced, and the search stops at a non-keyword value.)

26 mzlib/list

```
(require mzlib/list)
```

The `mzlib/list` library re-exports several functions from `scheme/base` and `scheme/list`:

```
cons?  
empty?  
empty  
foldl  
foldr  
remv  
remq  
remove  
remv*  
remq*  
remove*  
findf  
memf  
assf  
filter  
sort
```

```
(first v) → any/c  
  v : pair?  
(second v) → any/c  
  v : (and/c pair? ....)  
(third v) → any/c  
  v : (and/c pair? ....)  
(fourth v) → any/c  
  v : (and/c pair? ....)  
(fifth v) → any/c  
  v : (and/c pair? ....)  
(sixth v) → any/c  
  v : (and/c pair? ....)  
(seventh v) → any/c  
  v : (and/c pair? ....)  
(eighth v) → any/c  
  v : (and/c pair? ....)
```

Accesses the first, second, etc. element of “list” `v`. The argument need not actually be a list; it is inspected only as far as necessary to obtain an element (unlike the same-named functions from `scheme/list`, which do require the argument to be a list).

```
(rest v) → any/c  
v : pair?
```

The same as `cdr`.

```
(last-pair v) → pair?  
v : pair?
```

Returns the last pair in `v`, raising an error if `v` is not a pair (but `v` does not have to be a proper list).

```
(merge-sorted-lists lst1 lst2 less-than?) → list?  
lst1 : list?  
lst2 : list?  
less-than? : (any/c any/c . -> . any/c)
```

Merges the two sorted input lists, creating a new sorted list. The merged result is stable: equal items in both lists stay in the same order, and these in `lst1` precede `lst2`.

```
(mergesort lst less-than?) → list?  
lst : list?  
less-than? : (any/c any/c . -> . any/c)
```

The same as `sort`.

```
(quicksort lst less-than?) → list?  
lst : list?  
less-than? : (any/c any/c . -> . any/c)
```

The same as `sort`.

27 mzlib/match

```
(require mzlib/match)
```

The `mzlib/match` library provides a `match` form similar to that of `scheme/match`, but with an different (older and less extensible) syntax of patterns.

```
(match val-expr clause ...)
```

```
clause = [pat expr ...+]  
        | [pat (=> id) expr ...+]
```

See `match` from `scheme/match` for a description of matching. The grammar of `pat` for this `match` is as follows:

<code>pat</code>	<code>::= id</code>	match anything, bind identifier
	<code>-</code>	match anything
	<code>literal</code>	match literal
	<code>'datum</code>	match <code>equal?</code> datum
	<code>(lvp ...)</code>	match sequence of <code>lvps</code>
	<code>(lvp pat)</code>	match <code>lvps</code> consed onto a <code>pat</code>
	<code> #(lvp ...)</code>	match vector of <code>pat</code> s
	<code>#&pat</code>	match boxed <code>pat</code>
	<code>(\$ struct-id pat ...)</code>	match <code>struct-id</code> instance
	<code>(and pat ...)</code>	match when all <code>pat</code> s match
	<code>(or pat ...)</code>	match when any <code>pat</code> match
	<code>(not pat ...)</code>	match when no <code>pat</code> match
	<code>(= expr pat)</code>	match <code>(expr value)</code> to <code>pat</code>
	<code>(? pred-expr pat ...)</code>	match if <code>(expr value)</code> and <code>pat</code> s
	<code>'qp</code>	match quasipattern
<code>literal</code>	<code>::= #t</code>	match true
	<code>#f</code>	match false
	<code>string</code>	match <code>equal?</code> string
	<code>number</code>	match <code>equal?</code> number
	<code>character</code>	match <code>equal?</code> character
	<code>bytes</code>	match <code>equal?</code> byte string
	<code>keyword</code>	match <code>equal?</code> keyword
	<code>regexp</code>	match <code>equal?</code> regexp literal
	<code>pregexp</code>	match <code>equal?</code> pregexp literal
<code>lvp</code>	<code>::= pat</code>	greedily match <code>pat</code> instances
	<code>pat</code>	match <code>pat</code>
<code>ooo</code>	<code>::= ...</code>	zero or more; <code>...</code> is literal
	<code>---</code>	zero or more
	<code>..k</code>	<code>k</code> or more
	<code>__k</code>	<code>k</code> or more

<code>qp</code>	<code>::= literal</code>	match literal
	<code>id</code>	match <code>equal?</code> symbol
	<code>(qp ...)</code>	match sequences of <code>qps</code>
	<code>(qp qp)</code>	match sequence of <code>qps</code> consed onto a <code>qp</code>
	<code>(qp ... qp ooo)</code>	match <code>qps</code> consed onto a repeated <code>qp</code>
	<code> #(qp ...)</code>	match vector of <code>qps</code>
	<code> #&qp</code>	match boxed <code>qp</code>
	<code> ,pat</code>	match <code>pat</code>
	<code> ,@pat</code>	match <code>pat</code> , spliced

```

(match-lambda clause ...)
(match-lambda* clause ...)
(match-let ([pat expr] ...) body ...+)
(match-let* ([pat expr] ...) body ...+)
(match-letrec ([pat expr] ...) body ...+)
(match-define pat expr)

```

Analogous to the combined forms from `scheme/match`.

```

(define-match-expander id proc-expr)
(define-match-expander id proc-expr proc-expr)
(define-match-expander id proc-expr proc-expr proc-expr)
(match-equality-test) → (any/c any/c . -> . any)
(match-equality-test comp-proc) → void?
  comp-proc : (any/c any/c . -> . any)

```

Analogous to the form and parameter from `scheme/match`. The `define-match-expander` form, however, supports an extra `proc-expr` as the middle one: an expander for use with `match` from `mzlib/match`.

28 `mzlib/math`

```
(require mzlib/math)
```

Re-exports `scheme/math`, and also exports `e`.

`e` : `real?`

An approximation to Euler's constant: 2.718281828459045.

29 mzlib/md5

```
(require mzlib/md5)
```

Re-exports `file/md5`.

30 mzlib/os

```
(require mzlib/os)
```

```
(gethostname) → string?
```

Returns a string for the current machine's hostname (including its domain).

```
(getpid) → exact-integer?
```

Returns an integer identifying the current process within the operating system.

```
(truncate-file file [n-bytes]) → void?  
  file : path-string?  
  n-bytes : exact-nonnegative-integer? = 0
```

Truncates or extends the given *file* so that it is *n-bytes* long. If the file does not exist, or if the process does not have sufficient privilege to truncate the file, the `exn:fail` exception is raised.

WARNING: under Unix, the implementation assumes that the system's `ftruncate` function accepts a long long second argument.

31 mzlib/pconvert

```
(require mzlib/pconvert)
```

The `mzlib/pconvert` library defines routines for printing Scheme values as `evaluable` S-expressions, rather than `readable` S-expressions.

The `print-convert` procedure does not print values; rather, it converts a Scheme value into another Scheme value such that the new value pretty-prints as a Scheme expression that evaluates to the original value. For example, `(pretty-print (print-convert '(9 , (box 5) #(6 7))))` prints the literal expression `(list 9 (box 5) (vector 6 7))` to the current output port.

To install print converting into the read-eval-print loop, require `mzlib/pconvert` and call the procedure `install-converting-printer`.

In addition to `print-convert`, this library provides `print-convert`, `build-share`, `get-shared`, and `print-convert-expr`. The last three are used to convert sub-expressions of a larger expression (potentially with shared structure).

See also `prop:print-convert-constructor-name`.

```
(abbreviate-cons-as-list) → boolean?  
(abbreviate-cons-as-list abbreviate?) → void?  
  abbreviate? : any/c
```

A parameter that controls how lists are represented with constructor-style conversion. If the parameter's value is `#t`, lists are represented using `list`. Otherwise, lists are represented using `cons`. The initial value of the parameter is `#t`.

```
(booleans-as-true/false) → boolean?  
(booleans-as-true/false use-name?) → void?  
  use-name? : any/c
```

A parameter that controls how `#t` and `#f` are represented. If the parameter's value is `#t`, then `#t` is represented as `true` and `#f` is represented as `false`. The initial value of the parameter is `#t`.

```
(use-named/undefined-handler) → (any/c . -> . any/c)  
(use-named/undefined-handler use-handler) → void?  
  use-handler : (any/c . -> . any/c)
```

A parameter that controls how values that have inferred names are represented. The procedure is passed a value. If the procedure returns true, the procedure associated with

`named/undefined-handler` is invoked to render that value. Only values that have inferred names but are not defined at the top-level are used with this handler.

The initial value of the parameter is `(lambda (x) #f)`.

```
(named/undefined-handler) → (any/c . -> . any/c)
(named/undefined-handler use-handler) → void?
  use-handler : (any/c . -> . any/c)
```

Parameter for a procedure that controls how values that have inferred names are represented. The procedure is called only if `use-named/undefined-handler` returns true for some value. In that case, the procedure is passed that same value, and the result of the parameter is used as the representation for the value.

The initial value of the parameter is `(lambda (x) #f)`.

```
(build-share v) → ....
  v : any/c
```

Takes a value and computes sharing information used for representing the value as an expression. The return value is an opaque structure that can be passed back into `get-shared` or `print-convert-expr`.

```
(constructor-style-printing) → boolean?
(constructor-style-printing use-constructors?) → void?
  use-constructors? : any/c
```

Parameter that controls how values are represented after conversion. If this parameter's value is `#t`, then constructors are used; e.g., pair containing `1` and `2` is represented as `(cons 1 2)`. Otherwise, `quasiquote`-style syntax is used; e.g., the pair containing `1` and `2` is represented as `'(1 . 2)`. The initial value of the parameter is `#f`.

The constructor used for mutable pairs is `mcons`, unless `print-mpair-curly-braces` is set to `#f`, in which case `cons` and `list` are used. Similarly, when using `quasiquote` style and `print-mpair-curly-braces` is set to `#f`, mutable pair constructions are represented using `quote`, `quasiquote`, etc.

See also `quasi-read-style-printing` and `prop:print-convert-constructor-name`.

```
(current-build-share-hook)
→ (any/c (any/c . -> . void?)
        (any/c . -> . void?) . -> . any)
(current-build-share-hook hook) → void?
```

```
hook : (any/c (any/c . -> . void?)
         (any/c . -> . void?) . -> . any)
```

Parameter that sets a procedure used by `print-convert` and `build-share` to assemble sharing information. The procedure `hook` takes three arguments: a value `v`, a procedure `basic-share`, and a procedure `sub-share`; the return value is ignored. The `basic-share` procedure takes `v` and performs the built-in sharing analysis, while the `sub-share` procedure takes a component of `v` and analyzes it. Sharing information is accumulated as values are passed to `basic-share` and `sub-share`.

A `current-build-share-hook` procedure usually works together with a `current-print-convert-hook` procedure.

```
(current-build-share-name-hook)
→ (any/c . -> . (or/c symbol? false/c))
(current-build-share-name-hook hook) → void?
hook : (any/c . -> . (or/c symbol? false/c))
```

Parameter that sets a procedure used by `print-convert` and `build-share` to generate a new name for a shared value. The `hook` procedure takes a single value and returns a symbol for the value's name. If `hook` returns `#f`, a name is generated using the form “-*n*”, where *n* is an integer.

```
(current-print-convert-hook) → (any/c/ (any/c . -> . any/c)
                                       (any/c . -> . any/c))
(current-print-convert-hook hook) → void?
hook : (any/c/ (any/c . -> . any/c)
          (any/c . -> . any/c))
```

Parameter that sets a procedure used by `print-convert` and `print-convert-expr` to convert values. The procedure `hook` takes three arguments—a value `v`, a procedure `basic-convert`, and a procedure `sub-convert`—and returns the converted representation of `v`. The `basic-convert` procedure takes `v` and returns the default conversion, while the `sub-convert` procedure takes a component of `v` and returns its conversion.

A `current-print-convert-hook` procedure usually works together with a `current-build-share-hook` procedure.

```
(current-read-eval-convert-print-prompt) → string?
(current-read-eval-convert-print-prompt str) → void?
str : string?
```

Parameter that sets the prompt used by `install-converting-printer`. The initial value is “|- ”.

```
(get-shared share-info [cycles-only?])
→ (list-of (cons/c symbol? any/c))
  share-info : ....
  cycles-only? : any/c = #f
```

The `share-info` value must be a result from `build-share`. The procedure returns a list matching variables to shared values within the value passed to `build-share`.

The default value for `cycles-only?` is `#f`; if it is not `#f`, `get-shared` returns only information about cycles.

For example,

```
(get-shared (build-share (shared ([a (cons 1 b)]
                                  [b (cons 2 a)]))
                a)))
```

might return the list

```
'((-1- (cons 1 -2-)) (-2- (cons 2 -1-)))
```

```
(install-converting-printer) → void?
```

Sets the current print handler to print values using `print-convert`. The current read handler is also set to use the prompt returned by `current-read-eval-convert-print-prompt`.

```
(print-convert v [cycles-only?]) → any/c
  v : any/c
  cycles-only? : any/c = (show-sharing)
```

Converts the value `v`. If `cycles-only?` is not `#f`, then only circular objects are included in the output.

```
(print-convert-expr share-info
                    v
                    unroll-once?) → any/c
  share-info : ....
  v : any/c
  unroll-once? : any/c
```

Converts the value `v` using sharing information `share-info`, which was previously returned by `build-share` for a value containing `v`. If the most recent call to `get-shared` with

`share-info` requested information only for cycles, then `print-convert-expr` will only display sharing among values for cycles, rather than showing all value sharing.

The `unroll-once?` argument is used if `v` is a shared value in `share-info`. In this case, if `unroll-once?` is `#f`, then the return value will be a shared-value identifier; otherwise, the returned value shows the internal structure of `v` (using shared value identifiers within `v`'s immediate structure as appropriate).

```
(quasi-read-style-printing) → boolean?  
(quasi-read-style-printing on?) → void?  
  on? : any/c
```

Parameter that controls how vectors and boxes are represented after conversion when the value of `constructor-style-printing` is `#f`. If `quasi-read-style-printing` is set to `#f`, then boxes and vectors are unquoted and represented using constructors. For example, the list of a box containing the number 1 and a vector containing the number 1 is represented as `'(, (box 1) , (vector 1))`. If the parameter's value is `#t`, then `#&....` and `#(....)` are used, e.g., `'(#&1 #(1))`. The initial value of the parameter is `#t`.

```
(show-sharing) → boolean?  
(show-sharing show?) → void?  
  show? : any/c
```

Parameter that determines whether sub-value sharing is conserved (and shown) in the converted output by default. The initial value of the parameter is `#t`.

```
(whole/fractional-exact-numbers) → boolean?  
(whole/fractional-exact-numbers whole-frac?) → void?  
  whole-frac? : any/c
```

Parameter that controls how exact, non-integer numbers are converted when the numerator is greater than the denominator. If the parameter's value is `#t`, the number is converted to the form `(+ integer fraction)` (i.e., a list containing '+, an exact integer, and an exact rational less than 1 and greater than -1). The initial value of the parameter is `#f`.

32 mzlib/pconvert-prop

```
(require mzlib/pconvert-prop)
```

```
prop:print-convert-constructor-name : property?  
(print-convert-named-constructor? v) → any  
  v : any/c  
(print-convert-constructor-name v) → any  
  v : any/c
```

The `prop:print-convert-constructor-name` property can be given a symbol value for a structure type. In that case, for constructor-style print conversion via `print-convert`, instances of the structure are shown using the symbol as the constructor name. Otherwise, the constructor name is determined by prefixing `make-` onto the result of `object-name`.

The `print-convert-named-constructor?` predicate recognizes instances of structure types that have the `prop:print-convert-constructor-name` property, and `print-convert-constructor-name` extracts the property value.

33 mzlib/plt-match

```
(require mzlib/plt-match)
```

The `mzlib/plt-match` library mostly re-provides `scheme/match`.

```
(define-match-expander id proc-expr)  
(define-match-expander id proc-expr proc-expr)  
(define-match-expander id proc-expr proc-expr proc-expr)
```

The same as the form from `mzlib/match`.

34 mzlib/port

```
(require mzlib/port)
```

The `mzlib/port` library mostly re-provides `scheme/port`.

```
(strip-shell-command-start in) → void?  
in : input-port?
```

Reads and discards a leading `#!` in `in` (plus continuing lines if the line ends with a backslash). Since `#!` followed by a forward slash or space is a comment, this procedure is not needed before reading Scheme expressions.

35 mzlib/pregexp

```
(require mzlib/pregexp)
```

The `mzlib/pregexp` library provides wrappers around `regexp-match`, etc. that coerce string and byte-string arguments to `pregexp` matchers instead of `regexp` matchers.

The library also re-exports: `pregexp`, and it re-exports `regexp-quote` as `pregexp-quote`.

```
(pregexp-match pattern
               input
               [start-pos
               end-pos
               output-port])
→ (or/c (listof (or/c (cons (or/c string? bytes?)
                           (or/c string? bytes?))
                       false/c))
      false/c)
pattern : (or/c string? bytes? regexp? byte-regexp?)
input   : (or/c string? bytes? input-port?)
start-pos : nonnegative-exact-integer? = 0
end-pos  : (or/c nonnegative-exact-integer? false/c) = #f
output-port : (or/c output-port? false/c) = #f

(pregexp-match-positions pattern
                        input
                        [start-pos
                        end-pos
                        output-port])
→ (or/c (listof (or/c (cons nonnegative-exact-integer?
                            nonnegative-exact-integer?)
                    false/c))
      false/c)
pattern : (or/c string? bytes? regexp? byte-regexp?)
input   : (or/c string? bytes? input-port?)
start-pos : nonnegative-exact-integer? = 0
end-pos  : (or/c nonnegative-exact-integer? false/c) = #f
output-port : (or/c output-port? false/c) = #f

(pregexp-split pattern
               input
               [start-pos
               end-pos]) → (listof (or/c string? bytes?))
pattern : (or/c string? bytes? regexp? byte-regexp?)
input   : (or/c string? bytes? input-port?)
start-pos : nonnegative-exact-integer? = 0
end-pos  : (or/c nonnegative-exact-integer? false/c) = #f
```

```
(pregexp-replace pattern input insert) → (or/c string? bytes?)
  pattern : (or/c string? bytes? regexp? byte-regexp?)
  input   : (or/c string? bytes?)
  insert  : (or/c string? bytes?
            (string? . -> . string?)
            (bytes? . -> . bytes?))

(pregexp-replace* pattern input insert) → (or/c string? bytes?)
  pattern : (or/c string? bytes? regexp? byte-regexp?)
  input   : (or/c string? bytes?)
  insert  : (or/c string? bytes?
            (string? . -> . string?)
            (bytes? . -> . bytes?))
```

Like `regexp-match`, etc., but a string `pattern` argument is compiled via `pregexp`, and a byte string `pattern` argument is compiled via `byte-pregexp`.

36 mzlib/pretty

```
(require mzlib/pretty)
```

Re-exports `scheme/pretty`.

37 mzlib/process

`(require mzlib/process)`

Re-exports `scheme/system`.

38 mzlib/restart

```
(require mzlib/restart)
```

```
(restart-mzscheme init-argv  
                  adjust-flag-table  
                  argv  
                  init-namespace) → boolean?  
init-argv : (vectorof string?)  
adjust-flag-table : (any/c . -> . any/c)  
argv : (vectorof string?)  
init-namespace : (-> any)
```

See [scheme/sandbox](#) for a more general way to simulate running a new PLT Scheme process.

Simulates starting MzScheme with the vector of command-line strings *argv*. The *init-argv*, *adjust-flag-table*, and *init-namespace* arguments are used to modify the default settings for command-line flags, adjust the parsing of command-line flags, and customize the initial namespace, respectively.

The vector of strings *init-argv* is read first with the standard MzScheme command-line parsing. Flags that load files or evaluate expressions (e.g., *-f* and *-e*) are ignored, but flags that set MzScheme’s modes (e.g., *-c* or *-j*) effectively set the default mode before *argv* is parsed.

Before *argv* is parsed, the procedure *adjust-flag-table* is called with a command-line flag table as accepted by [parse-command-line](#). The return value must also be a table of command-line flags, and this table is used to parse *argv*. The intent is to allow *adjust-flag-table* to add or remove flags from the standard set.

After *argv* is parsed, a new thread and a namespace are created for the “restarted” MzScheme. (The new namespace is installed as the current namespace in the new thread.) In the new thread, restarting performs the following actions:

- The *init-namespace* procedure is called with no arguments. The return value is ignored.
- Expressions and files specified by *argv* are evaluated and loaded. If an error occurs, the remaining expressions and files are ignored, and the return value for [restart-mzscheme](#) is set to *#f*.
- The [read-eval-print-loop](#) procedure is called, unless a flag in *init-argv* or *argv* disables it. When [read-eval-print-loop](#) returns, the return value for [restart-mzscheme](#) is set to *#t*.

Before evaluating command-line arguments, an exit handler is installed that immediately returns from [restart-mzscheme](#) with the value supplied to the handler. This exit handler

remains in effect when `read-eval-print-loop` is called (unless a command-line argument changes it). If `restart-mzscheme` returns normally, the return value is determined as described above.

Note that an error in a command-line expression followed by `read-eval-print-loop` produces a `#t` result. This is consistent with MzScheme's stand-alone behavior.

39 `mzlib/runtime-path`

```
(require mzlib/runtime-path)
```

Re-exports `scheme/runtime-path`.

40 mzlib/sandbox

```
(require mzlib/sandbox)
```

The `mzlib/sandbox` library mostly re-exports `scheme/sandbox`, but it provides a slightly different `make-evaluator` function.

The library re-exports the following bindings:

```
sandbox-init-hook
sandbox-reader
sandbox-input
sandbox-output
sandbox-error-output
sandbox-propagate-breaks
sandbox-coverage-enabled
sandbox-namespace-specs
sandbox-override-collection-paths
sandbox-security-guard
sandbox-path-permissions
sandbox-network-guard
sandbox-make-inspector
sandbox-eval-limits
kill-evaluator
break-evaluator
set-eval-limits
put-input
get-output
get-error-output
get-uncovered-expressions
call-with-limits
with-limits
exn:fail:resource?
exn:fail:resource-resource
```

```
(make-evaluator language
  requires
    input-program ...) → (any/c . -> . any)
language : (or/c module-path?
  (one-of/c 'r5rs 'beginner 'beginner-abbr
    'intermediate 'intermediate-lambda 'advanced)
  (list/c (one-of/c 'special) symbol?)
  (list/c (one-of/c 'special) symbol?)
  (cons/c (one-of/c 'begin) list?))
```

```

requires : (or/c (cons/c 'begin list?)
                (listof (or/c module-path? path?)))
input-program : any/c
(make-evaluator module-decl) → (any/c . -> . any)
module-decl : (or/c syntax? pair?)

```

Like `make-evaluator` or `make-module-evaluator`, but with several differences:

- The `language` argument can be one of a fixed set of symbols: `'r5rs`, etc. They are converted by adding a `(list 'special)` wrapper.
- If `requires` starts with `'begin`, then each element in the remainder of the list is effectively evaluated as a prefix to the program. Otherwise, it corresponds to the `#:requires` argument of `make-evaluator`.
- For each of `language` and `requires` that starts with `'begin`, the expressions are inspected to find top-level `require` forms (using symbolic equality to detect `require`), and the required modules are added to the `#:allow` list for `make-evaluator`.

41 mzlib/sendevent

```
(require mzlib/sendevent)
```

The `mzlib/sendevent` library provides a `send-event` function that works only on Mac OS X, and only when running in MrEd (though the library can be loaded in MzScheme).

```
(send-event receiver-bytes
            event-class-bytes
            event-id-bytes
            [direct-arg-v
            argument-list]) → any/c
receiver-bytes : (lambda (s) (and (bytes? s)
                                  (= 4 (bytes-length s))))
event-class-bytes : (lambda (s) (and (bytes? s)
                                     (= 4 (bytes-length s))))
event-id-bytes : (lambda (s) (and (bytes? s)
                                  (= 4 (bytes-length s))))
direct-arg-v : any/c = (void)
argument-list : list? = null
```

Calls `send-event` `scheme/gui/base`, if available, otherwise raises `exn:fail:unsupported`.

42 mzlib/serialize

```
(require mzlib/serialize)
```

The `mzlib/serialize` library provides the same bindings as `scheme/serialize`, except that `define-serializable-struct` and `define-serializable-struct/versions` are based on the syntax of `define-struct` from `mzscheme`.

```
(define-serializable-struct id-maybe-super (field-id ...) maybe-inspector-expr)  
(define-serializable-struct/versions id-maybe-super vers-num (field-id ...)   
                                     (other-version-clause ...)   
                                     maybe-inspector-expr)
```

```
id-maybe-super = id  
                | (id super-id)
```

```
maybe-inspector-expr =  
                        | inspector-expr
```

```
other-version-clause = (other-vers make-proc-expr  
                           cycle-make-proc-expr)
```

Like `define-serializable-struct` and `define-serializable-struct/versions`, but with the syntax of closer to `define-struct` of `mzscheme`.

43 mzlib/shared

```
(require mzlib/shared)
```

Re-exports `scheme/shared`.

44 mzlib/string

```
(require mzlib/string)
```

The `mzlib/string` library re-exports several functions from `scheme/base`:

```
real->decimal-string  
regexp-quote  
regexp-replace-quote  
regexp-match*  
regexp-match-positions*  
regexp-match-peek-positions*  
regexp-split  
regexp-match-exact?
```

It also re-exports `regexp-try-match` as `regexp-match/fail-without-reading`.

```
(glob->regexp [str  
             hide-dots?  
             case-sensitive?  
             simple?]) → (or/c regexp? byte-regexp?)  
str : (or/c string bytes?) = ?  
hide-dots? : any/c = #t  
case-sensitive? : any/c  
              = (eq? (system-path-convention-type) 'unix)  
simple? : any/c = #f
```

Produces a regexp for an input “glob pattern” `str`. A glob pattern is one that matches `*` with any string, `?` with a single character, and character ranges are the same as in regexps (unless `simple?` is true). In addition, the resulting regexp does not match strings that begin with `.`, unless `str` begins with `.` or `hide-dots?` is `#f`. The resulting regexp can be used with string file names to check the glob pattern. If the glob pattern is provided as a byte string, the result is a byte regexp.

The `case-sensitive?` argument determines whether the resulting regexp is case-sensitive.

If `simple?` is true, then ranges with `[...]` in `str` are treated as literal character sequences.

```
(string-lowercase! str) → void?  
str : (and/c string? (not/c immutable?))
```

Destructively changes `str` to contain only lowercase characters.

```
(string-uppercase! str) → void?
```

```
str : (and/c string? (not/c immutable?))
```

Destructively changes *str* to contain only uppercase characters.

```
(eval-string str [err-handler]) → list?  
str : (or/c string? bytes?)  
err-handler : (or/c false/c          = #f  
               (any/c . -> . any/c)  
               (-> any/c))
```

Reads and evaluates S-expressions from *str*, returning results for all of the expressions in the string. If any expression produces multiple results, the results are spliced into the resulting list. If *str* contains only whitespace and comments, an empty list is returned, and if *str* contains multiple expressions, the result will be contain multiple values from all subexpressions.

The *err-handler* argument can be:

- `#f` (the default) which means that errors are not caught;
- a one-argument procedure, which will be used with an exception (when an error occurs) and its result will be returned
- a thunk, which will be used to produce a result.

```
(expr->string expr) → string?  
expr : any/c
```

Prints *expr* into a string and returns the string.

```
(read-from-string str [err-handler]) → any/c  
str : (or/c string? bytes?)  
err-handler : (or/c false/c          = #f  
               (any/c . -> . any/c)  
               (-> any/c))
```

Reads the first S-expression from *str* and returns it. The *err-handler* is as in `eval-string`.

```
(read-from-string-all str [err-handler]) → list?  
str : (or/c string? bytes?)
```

```
err-handler : (or/c false/c                = #f  
                (any/c . -> . any/c)  
                (-> any/c))
```

Reads all S-expressions from the string (or byte string) *str* and returns them in a list. The *err-handler* is as in *eval-string*.

45 mzlib/struct

```
(require mzlib/struct)
```

```
(copy-struct struct-id struct-expr  
            (accessor-id field-expr) ...)
```

“Functional update” for structure instances. The result of evaluating *struct-expr* must be an instance of the structure type named by *struct-id*. The result of the `copy-struct` expression is a fresh instance of *struct-id* with the same field values as the result of *struct-expr*, except that the value for the field accessed by each *accessor-id* is replaced by the result of *field-expr*.

The result of *struct-expr* might be an instance of a sub-type of *struct-id*, but the result of the `copy-struct` expression is an immediate instance of *struct-id*. If *struct-expr* does not produce an instance of *struct-id*, the `exn:fail:contract` exception is raised.

If any *accessor-id* is not bound to an accessor of *struct-id* (according to the expansion-time information associated with *struct-id*), or if the same *accessor-id* is used twice, then a syntax error is raised.

```
(define-struct/properties id (field-id ...)
                          ((prop-expr val-expr) ...)
                          maybe-inspector-expr)
```

```
maybe-inspector-expr =  
    | expr
```

Like `define-struct` from `mzscheme`, but properties can be attached to the structure type. Each *prop-expr* should produce a structure-type property value, and each *val-expr* produces the corresponding value for the property.

Examples:

```
> (define-struct/properties point (x y)
    ([prop:custom-write (lambda (p port write?)
                          (fprintf port "(~a, ~a)"
                                    (point-x p)
                                    (point-y p))))])
> (display (make-point 1 2))
(1, 2)
```

```
(make->vector struct-id)
```

Builds a function that accepts a structure type instance (matching *struct-id*) and provides a vector of the fields of the structure type instance.

46 `mzlib/stxparam`

```
(require mzlib/stxparam)
```

Re-exports `scheme/stxparam` and `scheme/stxparam-exptime` (both at phase level 0).

47 mzlib/surrogate

```
(require mzlib/surrogate)
```

Re-exports `scheme/surrogate`.

48 `mzlib/tar`

```
(require mzlib/tar)
```

Re-exports `file/tar`.

49 mzlib/thread

(require mzlib/thread)

```
(coroutine proc) → coroutine?  
proc : ((any/c . -> . void?) . -> . any/c)
```

Returns a coroutine object to encapsulate a thread that runs only when allowed. The `proc` procedure should accept one argument, and `proc` is run in the coroutine thread when `coroutine-run` is called. If `coroutine-run` returns due to a timeout, then the coroutine thread is suspended until a future call to `coroutine-run`. Thus, `proc` only executes during the dynamic extent of a `coroutine-run` call.

The argument to `proc` is a procedure that takes a boolean, and it can be used to disable suspends (in case `proc` has critical regions where it should not be suspended). A true value passed to the procedure enables suspends, and `#f` disables suspends. Initially, suspends are allowed.

```
(coroutine? v) → any  
v : any/c
```

Returns `#t` if `v` is a coroutine produced by `coroutine`, `#f` otherwise.

```
(coroutine-run until coroutine) → boolean?  
until : (or/c evt? real?)  
coroutine : coroutine?
```

Allows the thread associated with `coroutine` to execute for up as long as `until` milliseconds (of `until` is a real number) or `until` is ready (if `until` is an event). If `coroutine`'s procedure disables suspends, then the coroutine can run arbitrarily long until it re-enables suspends.

The `coroutine-run` procedure returns `#t` if `coroutine`'s procedure completes (or if it completed earlier), and the result is available via `coroutine-result`. The `coroutine-run` procedure returns `#f` if `coroutine`'s procedure does not complete before it is suspended after `timeout-secs`. If `coroutine`'s procedure raises an exception, then it is re-raised by `coroutine-run`.

```
(coroutine-result coroutine) → any  
coroutine : coroutine
```

Returns the result for `coroutine` if it has completed with a value (as opposed to an excep-

tion), #f otherwise.

```
(coroutine-kill coroutine) → void?  
  coroutine : coroutine?
```

Forcibly terminates the thread associated with *coroutine* if it is still running, leaving the coroutine result unchanged.

```
(consumer-thread f [init]) → thread? procedure?  
  f : procedure?  
  init : (-> any) = void
```

Returns two values: a thread descriptor for a new thread, and a procedure with the same arity as *f*.

When the returned procedure is applied, its arguments are queued to be passed on to *f*, and #<void> is immediately returned. The thread created by `consumer-thread` dequeues arguments and applies *f* to them, removing a new set of arguments from the queue only when the previous application of *f* has completed; if *f* escapes from a normal return (via an exception or a continuation), the *f*-applying thread terminates.

The *init* argument is a procedure of no arguments; if it is provided, *init* is called in the new thread immediately after the thread is created.

```
(run-server port-no  
           conn-proc  
           conn-timeout  
           [handler  
           listen  
           close  
           accept  
           accept/break]) → void?  
port-no : (integer-in 1 65535)  
conn-proc : (input-port? output-port? . -> . any)  
conn-timeout : (and/c real? (not/c negative?))  
handler : (exn? . -> . any/c) = void  
listen : ((integer-in 1 65535) (one-of/c 5) (one-of/c #t)  
          . -> . listener?)  
          = tcp-listen  
close : (listener? . -> . any) = tcp-close  
accept : (listener? . ->* . (input-port? output-port?))  
          = tcp-accept  
accept/break : (listener? . ->* . (input-port? output-port?))  
                = tcp-accept/enable-break
```

Executes a TCP server on the port indicated by `port-no`. When a connection is made by a client, `conn` is called with two values: an input port to receive from the client, and an output port to send to the client.

Each client connection is managed by a new custodian, and each call to `conn` occurs in a new thread (managed by the connection's custodian). If the thread executing `conn` terminates for any reason (e.g., `conn` returns), the connection's custodian is shut down. Consequently, `conn` need not close the ports provided to it. Breaks are enabled in the connection thread if breaks are enabled when `run-server` is called.

To facilitate capturing a continuation in one connection thread and invoking it in another, the parameterization of the `run-server` call is used for every call to `handler`. In this parameterization and for the connection's thread, the `current-custodian` parameter is assigned to the connection's custodian.

If `conn-timeout` is not `#f`, then it must be a non-negative number specifying the time in seconds that a connection thread is allowed to run before it is sent a break signal. Then, if the thread runs longer than `(* conn-timeout 2)` seconds, then the connection's custodian is shut down. If `conn-timeout` is `#f`, a connection thread can run indefinitely.

If `handler` is provided, it is passed exceptions related to connections (i.e., exceptions not caught by `conn-proc`, or exceptions that occur when trying to accept a connection). The default handler ignores the exception and returns `#<void>`.

The `run-server` function uses `listen`, `close`, `accept` and `accept/break` in the same way as it might use `tcp-listen`, `tcp-close`, `tcp-accept`, and `tcp-accept/enable-break` to accept connections. Provide alternate procedures to use an alternate communication protocol (such as SSL) or to supply optional arguments in the use of `tcp-listen`. The `listener?` part of the contract indicates that the procedures must all work on the same kind of listener value.

The `run-server` procedure loops to serve client connections, so it never returns. If a break occurs, the loop will cleanly shut down the server, but it will not terminate active connections.

50 `mzlib/trace`

```
(require mzlib/trace)
```

The `mzlib/trace` library mimics the tracing facility available in Chez Scheme.

```
(trace id ...)
```

Each *id* must be bound to a procedure in the environment of the `trace` expression. Each *id* is `set!`ed to a new procedure that traces procedure calls and returns by printing the arguments and results of the call. If multiple values are returned, each value is displayed starting on a separate line.

When traced procedures invoke each other, nested invocations are shown by printing a nesting prefix. If the nesting depth grows to ten and beyond, a number is printed to show the actual nesting depth.

The `trace` form can be used on an identifier that is already traced. In this case, assuming that the variable's value has not been changed, `trace` has no effect. If the variable has been changed to a different procedure, then a new trace is installed.

Tracing respects tail calls to preserve loops, but its effect may be visible through continuation marks. When a call to a traced procedure occurs in tail position with respect to a previous traced call, then the tailness of the call is preserved (and the result of the call is not printed for the tail call, because the same result will be printed for an enclosing call). Otherwise, however, the body of a traced procedure is not evaluated in tail position with respect to a call to the procedure.

The result of a `trace` expression is `#<void>`.

```
(untrace id ...)
```

Undoes the effects of the `trace` form for each *id*, `set!`ing each *id* back to the untraced procedure, but only if the current value of *id* is a traced procedure. If the current value of a *id* is not a procedure installed by `trace`, then the variable is not changed.

The result of an `untrace` expression is `#<void>`.

51 mzlib/traceld

```
(require mzlib/traceld)
```

The `mzlib/traceld` library does not provide any bindings. Instead, `mzlib/traceld` is required for its side-effects.

The `mzlib/traceld` library installs a new load handler (see [current-load](#)) and load-extension handler (see [current-load-extension](#)) to print information about the files that are loaded. These handlers chain to the current handlers to perform the actual loads. Trace output is printed to the port that is the current error port (see [current-error-port](#)) when the library is instantiated.

Before a file is loaded, the tracer prints the file name and “time” (as reported by the procedure [current-process-milliseconds](#)) when the load starts. Trace information for nested loads is printed with indentation. After the file is loaded, the file name is printed with the “time” that the load completed.

52 `mzlib/trait`

```
(require mzlib/trait)
```

Re-exports `scheme/trait`.

53 mzlib/transcr

```
(require mzlib/transcr)
```

The `transcript-on` and `transcript-off` procedures of `mzscheme` always raise `exn:fail:unsupported`. The `mzlib/transcr` library provides working versions of `transcript-on` and `transcript-off`.

```
(transcript-on filename) → any  
  filename : any/c  
(transcript-off) → any
```

Starts/stops recording a transcript at *filename*.

54 mzlib/unit

```
(require mzlib/unit)
```

The `mzlib/unit` library mostly re-provides `scheme/unit`, except for `struct` from `scheme/unit`.

```
(struct id (field-id ...) omit-decl ...)
```

```
omit-decl = -type  
           | -selectors  
           | -setters  
           | -constructor
```

A signature form like `struct` from `scheme/unit`, but with a different syntax for the options that limit exports.

55 `mzlib/unit-exptime`

```
(require mzlib/unit-exptime)
```

Re-exports `scheme/unit-exptime`.

56 `mzlib/unit200`

```
(require mzlib/unit200)
```

The `mzlib/unit200` library provides an old implementation of units. See archived version 360 documentation on the `"unit.ss"` library of the `"mzlib"` collection for information about this library.

57 `mzlib/unitsig200`

```
(require mzlib/unitsig200)
```

The `mzlib/unit200` library provides an old implementation of units. See archived version 360 documentation on the "unitsig.ss" library of the "mzlib" collection for information about this library.

58 mzlib/zip

`(require mzlib/zip)`

Re-exports `file/zip`.

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